

# Solutions - Homework 1

(Due date: Sep. 26<sup>th</sup> (008), Sep. 27<sup>th</sup> (010) @ 11:59 pm)

Presentation and clarity are very important! Show your procedure!

## PROBLEM 1 (28 PTS)

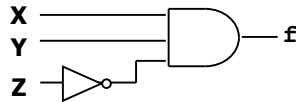
- a) Simplify the following functions using ONLY Boolean Algebra Theorems. For each resulting simplified function, sketch the logic circuit using AND, OR, XOR, and NOT gates. (15 pts)

$$\checkmark f = \overline{y(z + \bar{x})} + \overline{y\bar{x}}$$

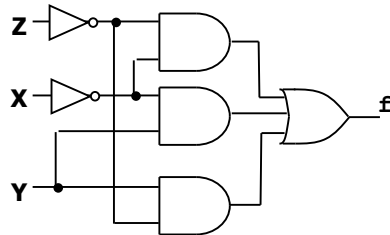
$$\checkmark f = \prod(M_1, M_4, M_5, M_7)$$

$$\checkmark f(A, B, C) = \overline{ABC} + \overline{(C \oplus A)B}$$

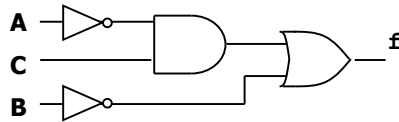
$$\checkmark f(x, y, z) = \overline{y(z + \bar{x})} + \overline{y\bar{x}} = \overline{YZ + Y\bar{X} + \bar{X} + \bar{Y}} = \overline{YZ + \bar{X} + \bar{Y}} = \overline{\bar{X} + (\bar{Y} + Y)(\bar{Y} + Z)} = \overline{\bar{X} + \bar{Y} + Z} = XY\bar{Z}$$



$$\checkmark f(X, Y, X) = \prod(M_1, M_4, M_5, M_7) = \sum(m_0, m_2, m_3, m_6) = \bar{X}\bar{Y}\bar{Z} + \bar{X}Y\bar{Z} + \bar{X}YZ + XY\bar{Z} = \bar{X}\bar{Z} + \bar{X}YZ + XY\bar{Z} \\ = \bar{X}(\bar{Z} + YZ) + XY\bar{Z} = \bar{X}(\bar{Z} + Y) + XY\bar{Z} = \bar{X}\bar{Z} + \bar{X}Y + XY\bar{Z} = \bar{X}Y + \bar{Z}(\bar{X} + XY) = \bar{X}Y + \bar{Z}(\bar{X} + Y) \\ = \bar{X}Y + \bar{Z}\bar{X} + \bar{Z}Y$$



$$\checkmark f(A, B, C) = \overline{ABC} + \overline{(A \oplus C)B} = \overline{ABC} + \overline{(AC + \bar{A}\bar{C})B} = \overline{AB(\bar{C} + C)} + \overline{\bar{A}\bar{C}B} = \overline{AB} + \overline{\bar{A}\bar{C}B} = \overline{B(A + \bar{A}\bar{C})} \\ = \overline{B(A + \bar{C})} = \bar{B} + \overline{(A + \bar{C})} = \bar{B} + \bar{A}C$$



- b) Using Boolean Algebra Theorems, prove that:  $y(x \oplus z) = (yx) \oplus (yz)$  (5 pts)

$$y(x \oplus z) = y(\bar{x}z + x\bar{z}) = y\bar{x}z + yx\bar{z}$$

$$(yx) \oplus (yz) = \bar{y}\bar{x}yz + yx\bar{y}z = (\bar{y} + \bar{x})yz + yx(\bar{y} + \bar{z}) = \bar{y}yz + \bar{x}yz + yx\bar{y} + yx\bar{z} = \bar{x}yz + yx\bar{z}$$

$$\therefore y(x \oplus z) = (yx) \oplus (yz)$$

c) For the following Truth table with two outputs: (8 pts)

- Provide the Boolean functions using the Canonical Sum of Products (SOP), and Product of Sums (POS). (4 pts)
- Express the Boolean functions using the minterms and maxterms representations.
- Sketch the logic circuits as Canonical Sum of Products and Product of Sums. (3 pts)

x	y	z	f <sub>1</sub>	f <sub>2</sub>
0	0	0	1	0
0	0	1	0	1
0	1	0	0	1
0	1	1	0	0
1	0	0	0	1
1	0	1	1	1
1	1	0	0	0
1	1	1	1	1

**Sum of Products**

$$f_1 = \bar{x}\bar{y}\bar{z} + x\bar{y}z + xyz$$

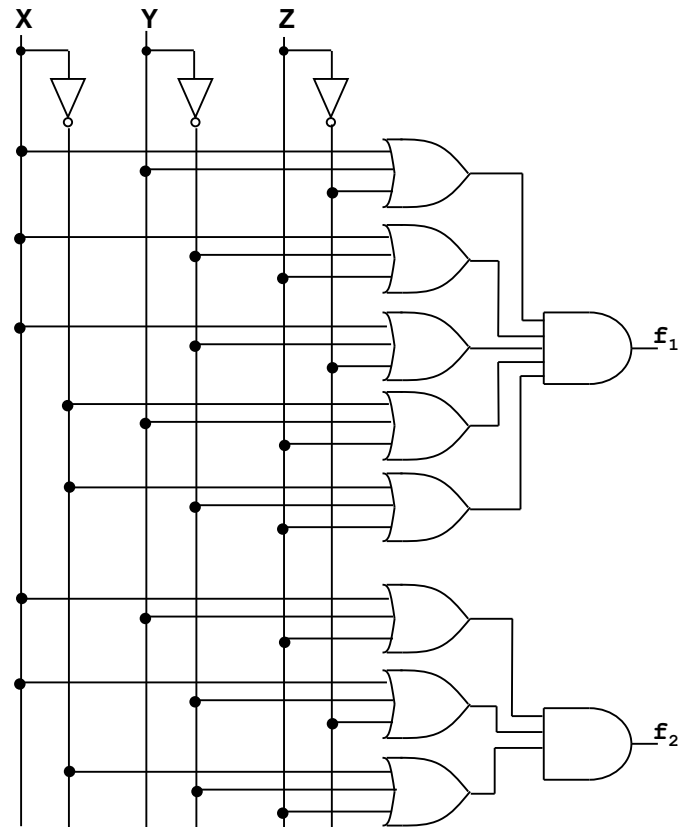
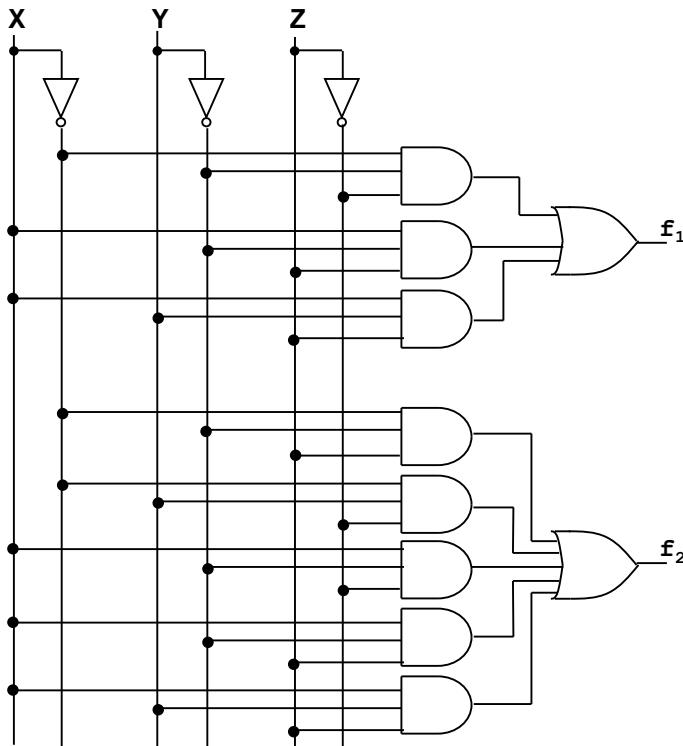
$$f_2 = \bar{x}\bar{y}z + \bar{x}y\bar{z} + x\bar{y}\bar{z} + x\bar{y}z + xyz$$

**Product of Sums**

$$f_1 = (x + y + \bar{z})(x + \bar{y} + z)(\bar{x} + y + z)(\bar{x} + \bar{y} + z)$$

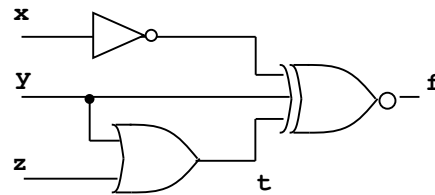
$$f_2 = (x + y + z)(x + \bar{y} + \bar{z})(\bar{x} + \bar{y} + z)$$

**Minterms and maxterms:**  $f_1 = \sum(m_0, m_5, m_7) = \prod(M_1, M_2, M_3, M_4, M_6)$ .  
 $f_2 = \sum(m_1, m_2, m_4, m_5, m_7) = \prod(M_0, M_3, M_6)$ .



## PROBLEM 2 (26 PTS)

- a) Construct the truth table describing the output of the following circuit and write the simplified Boolean equation (6 pts).  
Note that  $a \oplus b \oplus c = (a \oplus b) \oplus c = a \oplus (b \oplus c) = b \oplus (a \oplus c)$



x	y	z	t	f
0	0	0	0	0
0	0	1	1	1
0	1	0	1	0
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	1
1	1	1	1	1

$$f = \overline{x \oplus y \oplus (y + z)} = \overline{x \oplus (\bar{y}z)} = \bar{x}\bar{y}z + x\bar{y}z = \bar{x}\bar{y}z + xy + x\bar{z}$$

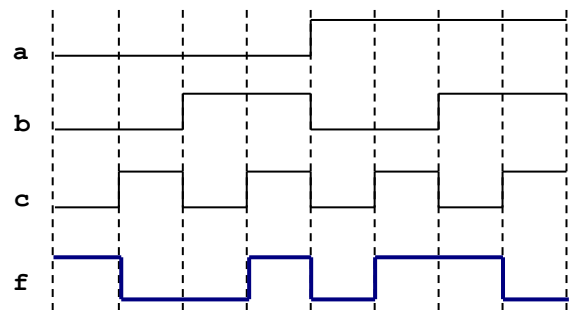
- b) The following is the timing diagram of a logic circuit with 3 inputs. Sketch the logic circuit that generates this waveform. Then, complete the VHDL code (using VHDL signals is optional). (8 pts)

```
library ieee;
use ieee.std_logic_1164.all;

entity circ is
  port ( a, b, c: in std_logic;
        f: out std_logic);
end circ;

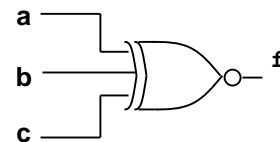
architecture st of circ is

begin
  f <= not (a xor b xor c);
end st;
```



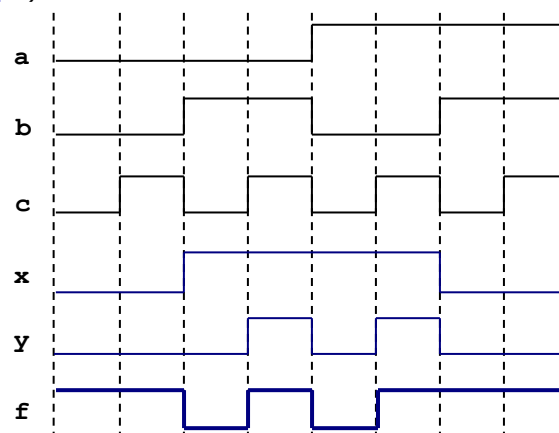
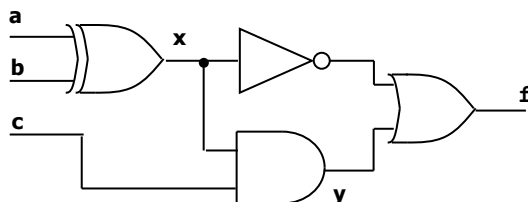
a	b	c	f
0	0	0	1
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	0

	ab	00	01	11	10
c	0	1	0	1	0
	1	0	1	0	1



$$f = \bar{a}\bar{b}\bar{c} + \bar{a}b\bar{c} + a\bar{b}\bar{c} + a\bar{b}c = \bar{a}(\bar{b}\bar{c} + b\bar{c}) + a(\bar{b}\bar{c} + \bar{b}c) = \bar{a}(\bar{b} \oplus c) + a(\bar{b} \oplus c) = \bar{a} \oplus (\bar{b} \oplus c) = \bar{a} \oplus b \oplus c$$

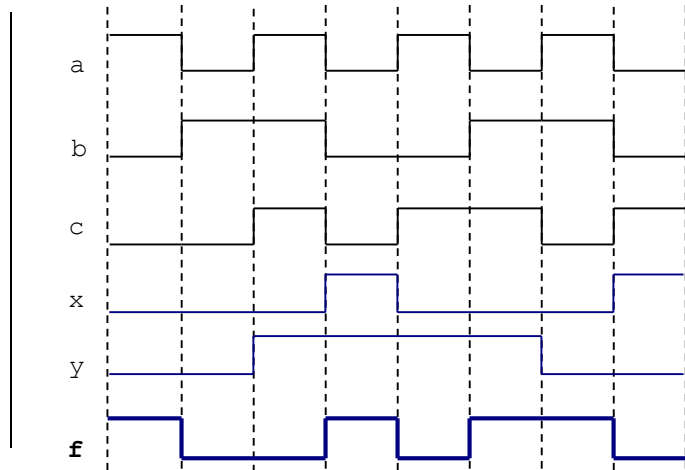
- c) Complete the timing diagram of the following circuit: (5 pts)



```
library ieee;
use ieee.std_logic_1164.all;

entity circ is
    port ( a, b, c: in std_logic;
           f: out std_logic);
end circ;

architecture st of circ is
    signal x, y: std_logic;
begin
    f <= (not y) xnor a;
    x <= a nor b;
    y <= x xor c;
end st;
```



- Complete the truth table for a circuit with 4 inputs  $x, y, z, w$  that activates an output ( $f = 1$ ) when the number of 1's in the inputs is equal than the number of 0's. For example: If  $xyzw = 1001 \rightarrow f = 1$ . If  $xyzw = 1011 \rightarrow f = 0$ .
- Provide the Boolean equation for the output  $f$  using the minterms representation.
- Sketch the logic circuit.

x	y	z	w	f
0	0	0	0	0
0	0	0	1	0
0	0	1	0	0
0	0	1	1	1
0	1	0	0	0
0	1	0	1	1
0	1	1	0	1
0	1	1	1	0
1	0	0	0	0
1	0	0	1	1
1	0	1	0	1
1	0	1	1	0
1	1	0	0	1
1	1	0	1	0
1	1	1	0	0
1	1	1	1	0

Karnaugh map for the function  $F(x, y, z, w)$ :

xy \ zw	00	01	11	10
00	0	0	1	0
01	0	1	0	1
11	1	0	0	0
10	0	1	0	1

Prime implicants (PIs) are circled in red:

- $\overline{x}\overline{y}\overline{z}\overline{w}$  (top-left 1)
- $\overline{x}yz\overline{w}$  (top-right 1)
- $\overline{x}yz\overline{w}$  (bottom-left 1)
- $\overline{x}yz\overline{w}$  (bottom-right 1)

The function is the sum of these prime implicants:

$$F(x, y, z, w) = \overline{x}\overline{y}\overline{z}\overline{w} + \overline{x}yz\overline{w} + \overline{x}yz\overline{w} + \overline{x}yz\overline{w}$$

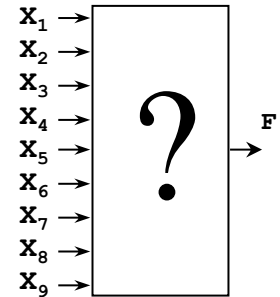
```

graph LR
    x((x)) --> bx(( ))
    x --> nx[NOT]
    y((y)) --> by(( ))
    y --> ny[NOT]
    z((z)) --> bz(( ))
    z --> nz[NOT]
    w((w)) --> bw(( ))
    w --> nw[NOT]
    
    AND1[AND]
    AND2[AND]
    AND3[AND]
    OR1[OR]
    
    bx --> AND1
    ny --> AND1
    nz --> AND1
    
    by --> AND2
    nx --> AND2
    nw --> AND2
    
    bz --> AND3
    bw --> AND3
    ny --> AND3
    
    AND1 --> OR1
    AND2 --> OR1
    AND3 --> OR1
    
    OR1 --> f((f))
  
```

### PROBLEM 4 (11 PTS)

- Tic-tac-toe game (3-by-3 grid of squares). The players alternate turns. Each player chooses a square and places a mark in a square (one player uses **x** and the other **o**). The first player with three marks in a row, column, or diagonal wins the game.
- Design a digital circuit for an electronic tic-tac-toe that indicates the presence of a winning pattern for player **x**. The circuit has 9 inputs ( $x_1$  to  $x_9$ ) and an output **F**.
  - Inputs  $x_1$  to  $x_9$ : A value of '1' indicates that the player marked the corresponding position with an **x**. A value of '0' indicates that the other player marked that position.
  - $F = 1$  if a winning pattern is present and  $F = 0$  otherwise.
  - Example: if  $x_1=1, x_2=0, x_3=1, x_4=0, x_5=1, x_6=0, x_7=1, x_8=0, x_9=1$ , then  $F=1$ .
- Provide the Boolean expression for **F**. The 9 inputs ( $x_1$  to  $x_9$ ) are arranged in the following pattern:  
Tip: Note that if there are three 1's in a winning pattern, the value of the remaining 6 positions is irrelevant.

$x_1$	$x_2$	$x_3$
$x_4$	$x_5$	$x_6$
$x_7$	$x_8$	$x_9$



$f = 1$  when there are three 1's in a line. Note that there can be more than three 1's before we find three 1's in a line. So, we make  $f = 1$  when a winning pattern is present regardless of the values of the remaining variables:

$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$	$x_7$	$x_8$	$x_9$	$f$
1	1	1	X	X	X	X	X	X	1
X	X	X	1	1	1	X	X	X	1
X	X	X	X	X	X	1	1	1	1
1	X	X	1	X	X	1	X	X	1
X	1	X	X	1	X	X	1	X	1
X	X	1	X	X	1	X	X	1	1
1	X	X	X	1	X	X	X	1	1
X	X	1	X	1	X	1	X	X	1
All remaining cases									0

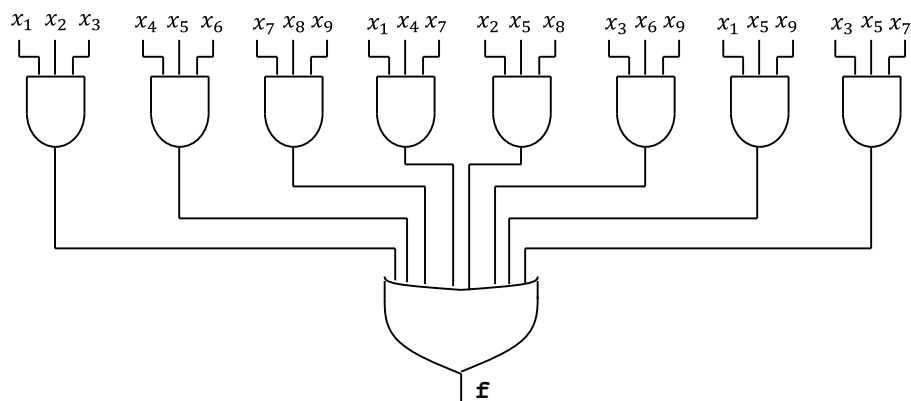
First row:  $x_1x_2x_3$ . (all permutations for  $x_4x_5x_6x_7x_8x_9$ ) =  $x_1x_2x_3 \cdot h(x_4, x_5, x_6, x_7, x_8, x_9)$

Where:  $h(x_4, x_5, x_6, x_7, x_8, x_9) = \sum m(0,1, \dots, 63) = 1$  (the sum of all the minterms of a 6-variable function is always 1).  
Thus, Then first row is just:  $x_1x_2x_3$

The same technique is applied for every row.

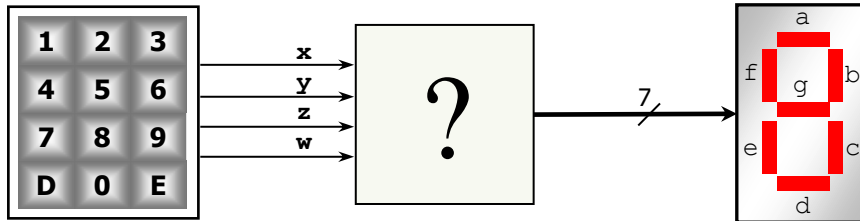
$$\rightarrow f = x_1x_2x_3 + x_4x_5x_6 + x_7x_8x_9 + x_1x_4x_7 + x_2x_5x_8 + x_3x_6x_9 + x_1x_5x_9 + x_3x_5x_7$$

- Sketch the logical circuit resulting from the Boolean equation for **F**.

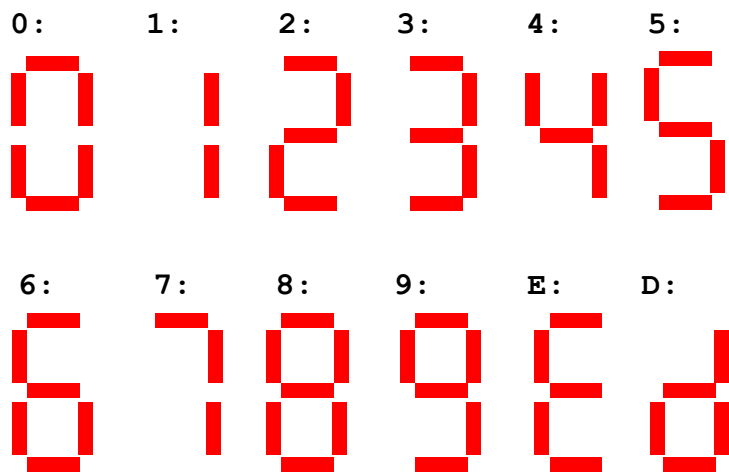


### PROBLEM 5 (25 PTS)

- A numeric keypad produces a 4-bit code as shown below. We want to design a logic circuit that converts each 4-bit code to a 7-segment code, where each segment is an LED. The LEDs are lit with a logical '0' (negative logic). The inputs are active high (or in positive logic).
- Complete the truth table for each output ( $a, b, c, d, e, f, g$ ). (4 pts)
- Provide the simplified expression for each output ( $a, b, c, d, e, f, g$ ). Use Karnaugh maps for  $a, b, c, d, e$  and the Quine-McCluskey algorithm for  $f, g$ . Note that it is safe to assume that the codes 1100 to 1111 will not be produced by the keypad.



Value	X	Y	Z	W	a	b	c	d	e	f	g
0	0	0	0	0							
1	0	0	0	1							
2	0	0	1	0							
3	0	0	1	1							
4	0	1	0	0							
5	0	1	0	1							
6	0	1	1	0							
7	0	1	1	1							
8	1	0	0	0							
9	1	0	0	1	0	0	0	0	1	0	0
E	1	0	1	0							
D	1	0	1	1							
	1	1	0	0							
	1	1	0	1							
	1	1	1	0							
	1	1	1	1							



Value	X	Y	Z	W	a	b	c	d	e	f	g
0	0	0	0	0	0	0	0	0	0	0	1
1	0	0	0	1	1	0	0	1	1	1	1
2	0	0	1	0	0	0	1	0	0	1	0
3	0	0	1	1	0	0	0	0	1	1	0
4	0	1	0	0	1	0	0	1	1	0	0
5	0	1	0	1	0	1	0	0	1	0	0
6	0	1	1	0	0	1	0	0	0	0	0
7	0	1	1	1	0	0	0	1	1	1	1
8	1	0	0	0	0	0	0	0	0	0	0
9	1	0	0	1	0	0	0	0	1	0	0
E	1	0	1	0	0	1	1	0	0	0	0
D	1	0	1	1	1	0	0	0	0	1	0
	1	1	0	0	X	X	X	X	X	X	X
	1	1	0	1	X	X	X	X	X	X	X
	1	1	1	0	X	X	X	X	X	X	X
	1	1	1	1	X	X	X	X	X	X	X

$$a = \bar{x}\bar{y}\bar{z}w + y\bar{z}\bar{w} + xzw$$

$$b = y\bar{z}w + yz\bar{w} + xz\bar{w}$$

$$c = \bar{y}z\bar{w}$$

$$d = \bar{x}\bar{y}\bar{z}w + y\bar{z}\bar{w} + yzw$$

$$e = y\bar{z} + \bar{z}w + \bar{x}w$$

a

xy \ zw	00	01	11	10
00	0	1	X	0
01	1	0	X	0
11	0	0	X	1
10	0	0	X	0

b

xy \ zw	00	01	11	10
00	0	0	X	0
01	0	1	X	0
11	0	0	X	0
10	0	1	X	1

c

xy \ zw	00	01	11	10
00	0	0	X	0
01	0	0	X	0
11	0	0	X	0
10	1	0	X	1

d

xy \ zw	00	01	11	10
00	0	1	X	0
01	1	0	X	0
11	0	1	X	0
10	0	0	X	0

e

xy \ zw	00	01	11	10
00	0	1	X	0
01	1	1	X	1
11	1	1	X	0
10	0	0	X	0

$$f = \sum m(1,2,3,7,11) + \sum d(12,13,14,15)$$

Number of ones	4-literal implicants	3-literal implicants	2-literal implicants	1-literal implicants
1	$m_1 = 0001$ ✓ $m_2 = 0010$ ✓	$m_{1,3} = 00-1$ $m_{2,3} = 001-$		We can't combine any further, so we stop here
2	$m_3 = 0011$ ✓ $m_{12} = 1100$ ✓	$m_{3,7} = 0-11$ ✓ $m_{3,11} = -011$ ✓ $m_{12,13} = 110-$ ✓ $m_{12,14} = 11-0$ ✓	$m_{3,7,11,15} = --11$ <del><math>m_{3,11,7,15} = --11</math></del> $m_{14,15,12,13} = 11--$ <del><math>m_{12,14,13,15} = 11--</math></del>	
3	$m_7 = 0111$ ✓ $m_{11} = 1011$ ✓ $m_{13} = 1101$ ✓ $m_{14} = 1110$ ✓	$m_{7,15} = -111$ ✓ $m_{11,15} = 1-11$ ✓ $m_{13,15} = 11-1$ ✓ $m_{14,15} = 111-$ ✓		
4	$m_{15} = 1111$ ✓			

$$f = \bar{x}\bar{y}w + \bar{x}\bar{y}z + zw + xy$$

Prime Implicants		Minterms				
		1	2	3	7	11
$m_{1,3}$	$\bar{x}\bar{y}w$	X		X		
$m_{2,3}$	$\bar{x}\bar{y}z$		X	X		
$m_{3,7,11,15}$	$zw$			X	X	X
$m_{14,15,12,13}$	$xy$					

$$\rightarrow f = \bar{x}\bar{y}w + \bar{x}\bar{y}z + zw$$

$$g = \sum m(0,1,7) + \sum d(12,13,14,15).$$

Number of ones	4-literal implicants	3-literal implicants	2-literal implicants	1-literal implicants
0	$m_0 = 0000$ ✓	$m_{0,1} = 000-$		
1	$m_1 = 0001$ ✓			
2	$m_{12} = 1100$ ✓	$m_{12,13} = 110-$ ✓ $m_{12,14} = 11-0$ ✓	$m_{12,13,14,15} = 11--$ <del><math>m_{12,14,13,15} = 11--</math></del>	We can't combine any further, so we stop here
3	$m_7 = 0111$ ✓ $m_{13} = 1101$ ✓ $m_{14} = 1110$ ✓	$m_{7,15} = -111$ ✓ $m_{13,15} = 11-1$ ✓ $m_{14,15} = 111-$ ✓		
4	$m_{15} = 1111$ ✓			

$$g = \bar{x}\bar{y}\bar{z} + yzw + xy$$

Prime Implicants		Minterms		
		0	1	7
$m_{0,1}$	$\bar{x}\bar{y}\bar{z}$	<b>X</b>	X	
$m_{7,15}$	$yzw$			<b>X</b>
$m_{12,13,14,15}$	$xy$			

$$\rightarrow g = \bar{x}\bar{y}\bar{z} + yzw$$